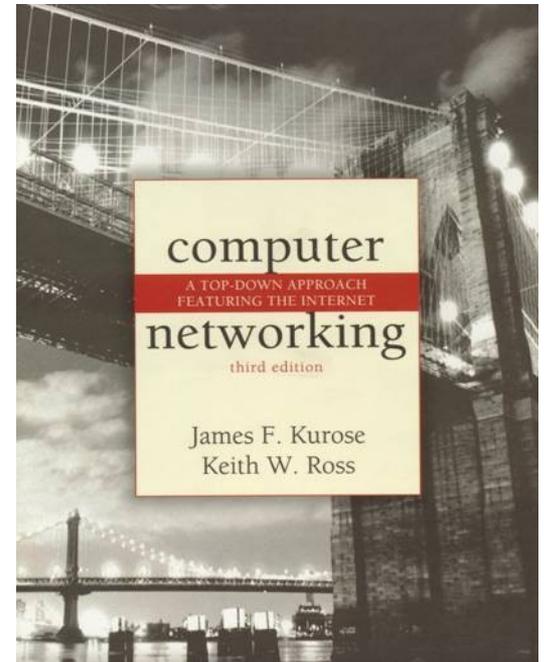


# Wireless Networks



*Computer Networking:  
A Top Down Approach  
Featuring the Internet,*

*3<sup>rd</sup> edition.*

*Jim Kurose, Keith Ross  
Addison-Wesley, July  
2004.*

# Wireless Networks

## Background:

- ❑ # wireless (mobile) phone subscribers now exceeds # wired phone subscribers!
- ❑ computer nets: laptops, palmtops, PDAs, Internet-enabled phone promise anytime untethered Internet access
- ❑ two important (but different) challenges
  - communication over wireless link
  - handling mobile user who changes point of attachment to network

# Outline

## Introduction

### Wireless

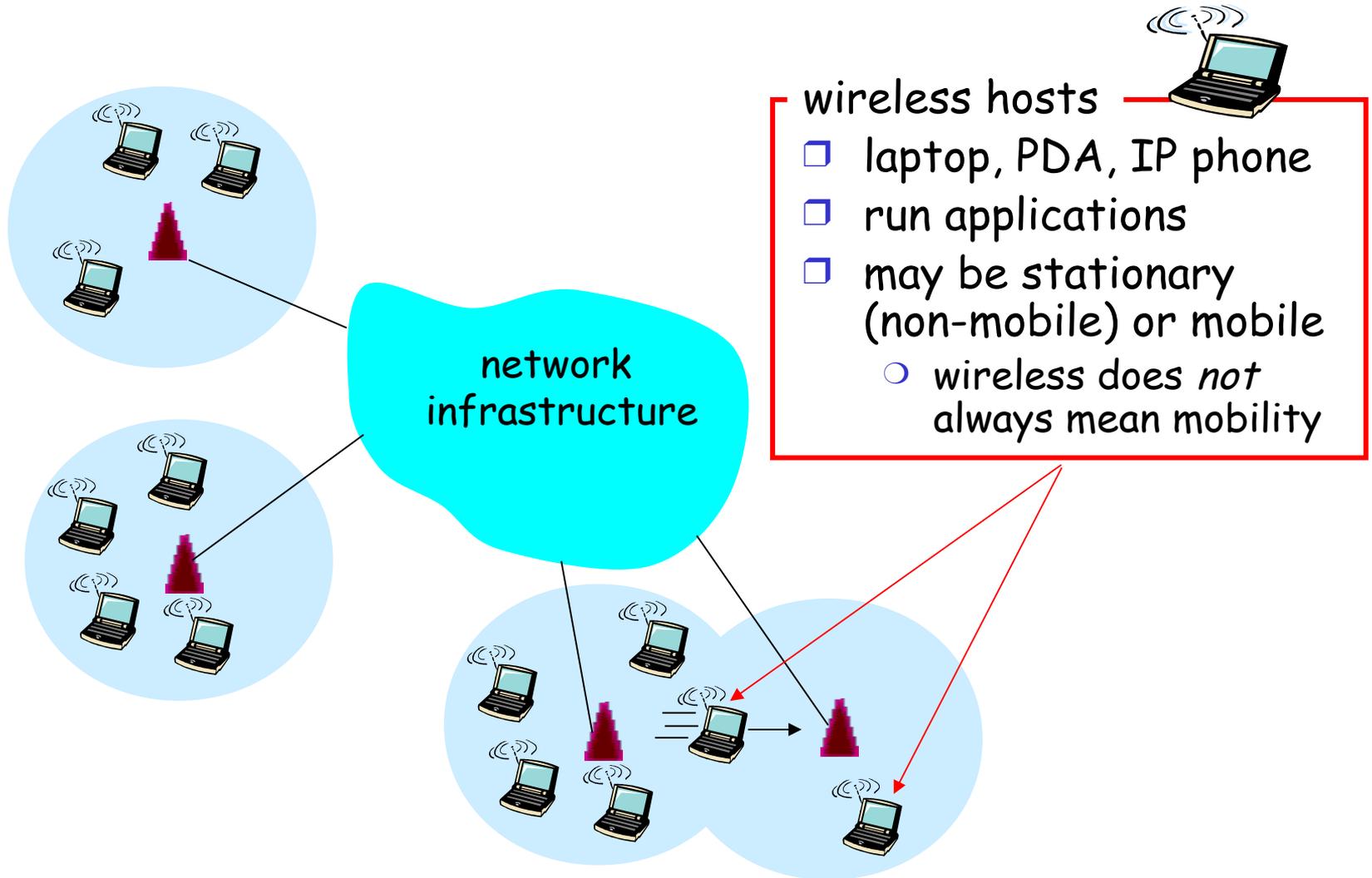
- Wireless links, characteristics
- IEEE 802.11 wireless LANs ("wi-fi")

### Mobility

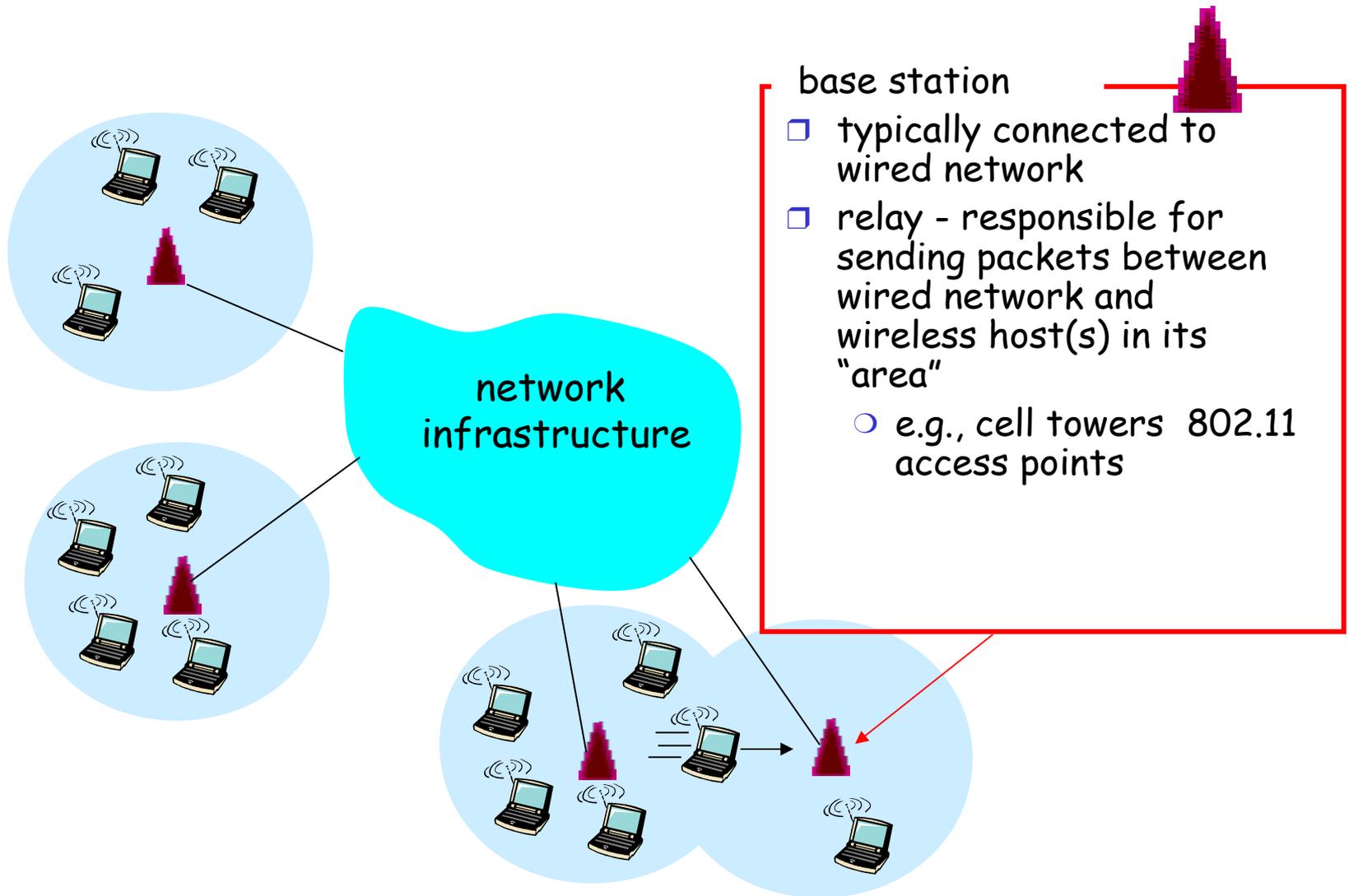
- Principles: addressing and routing to mobile users
- Mobile IP
- Mobility and higher-layer protocols

### Summary

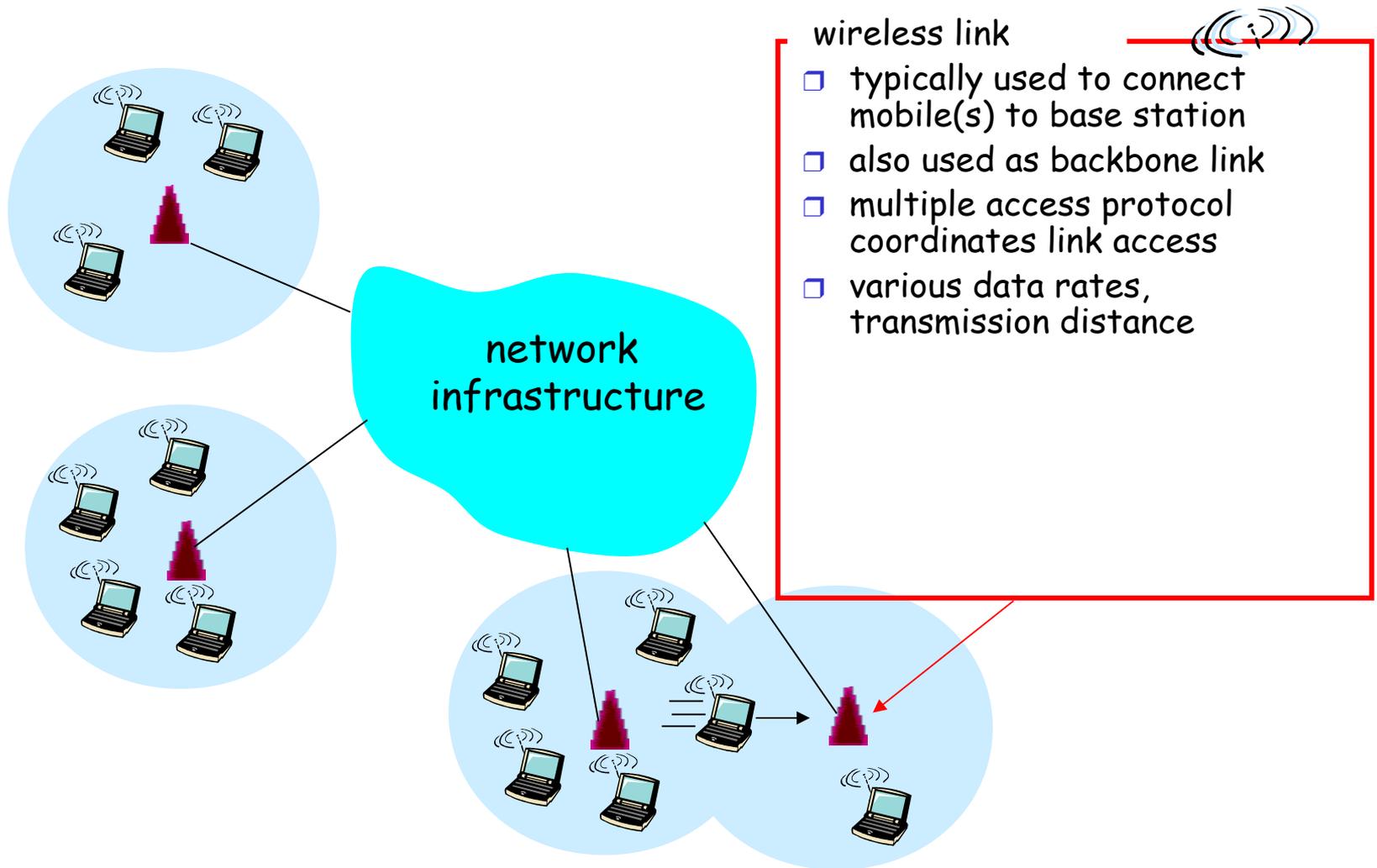
# Elements of a wireless network



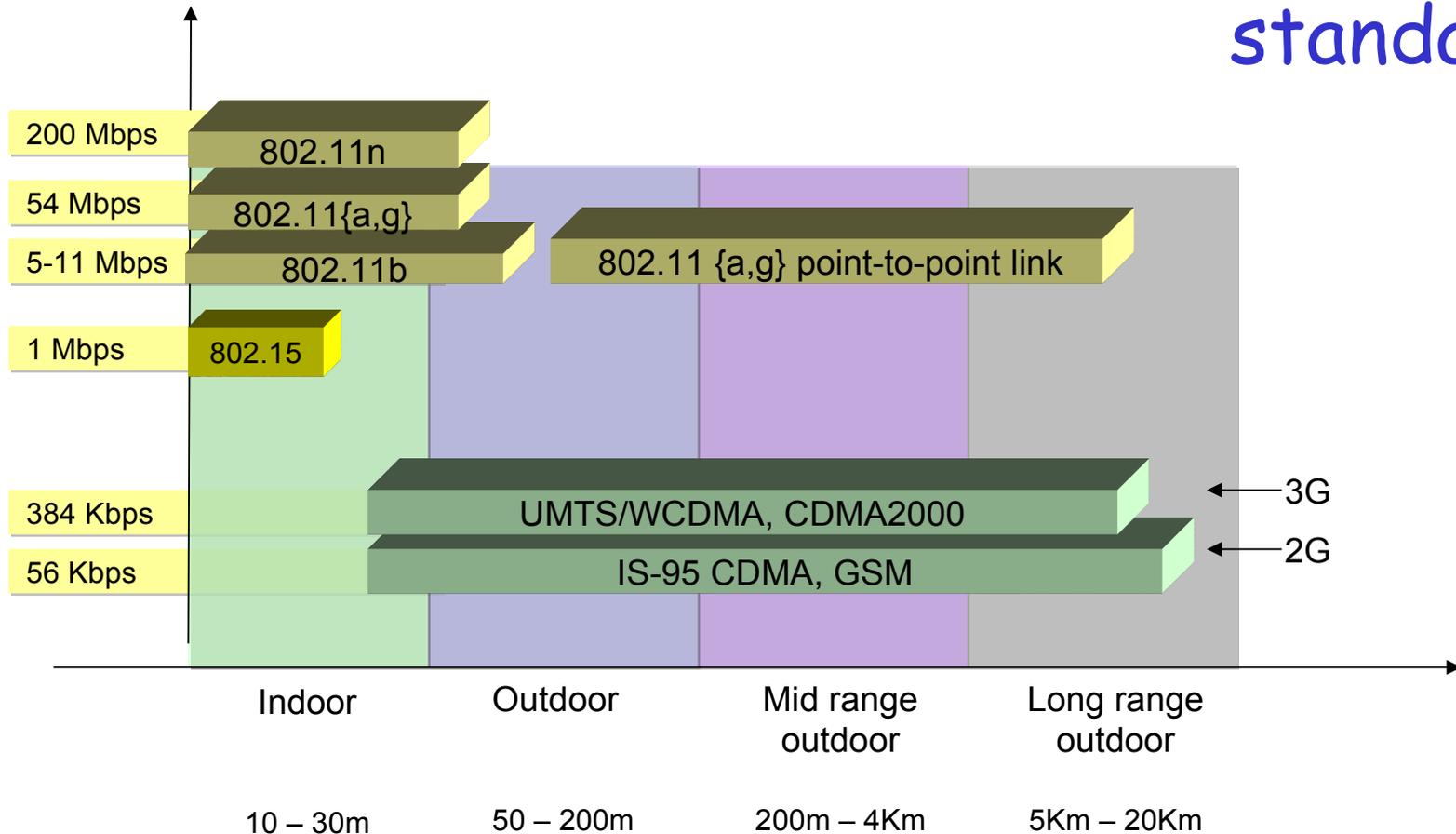
# Elements of a wireless network



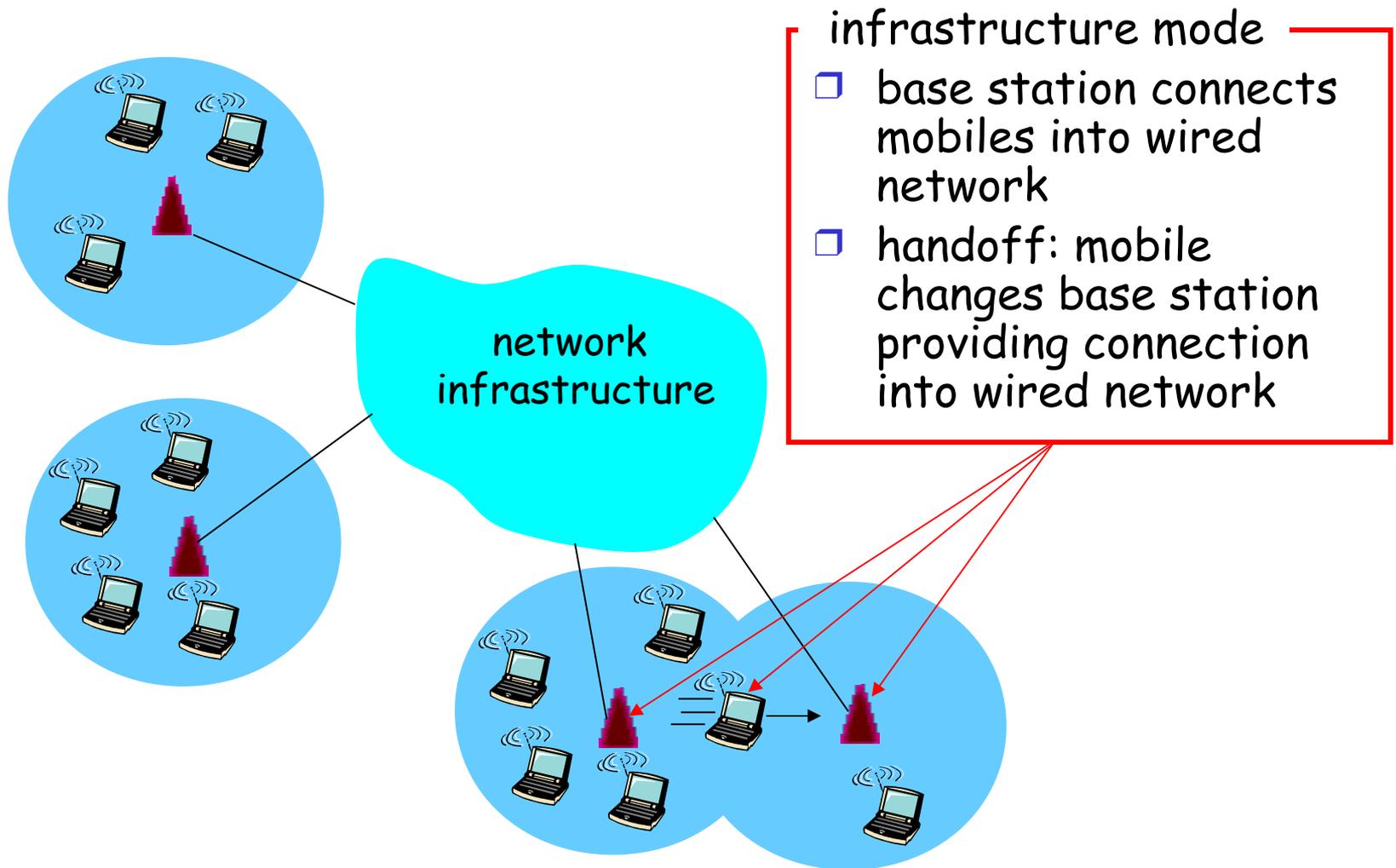
# Elements of a wireless network



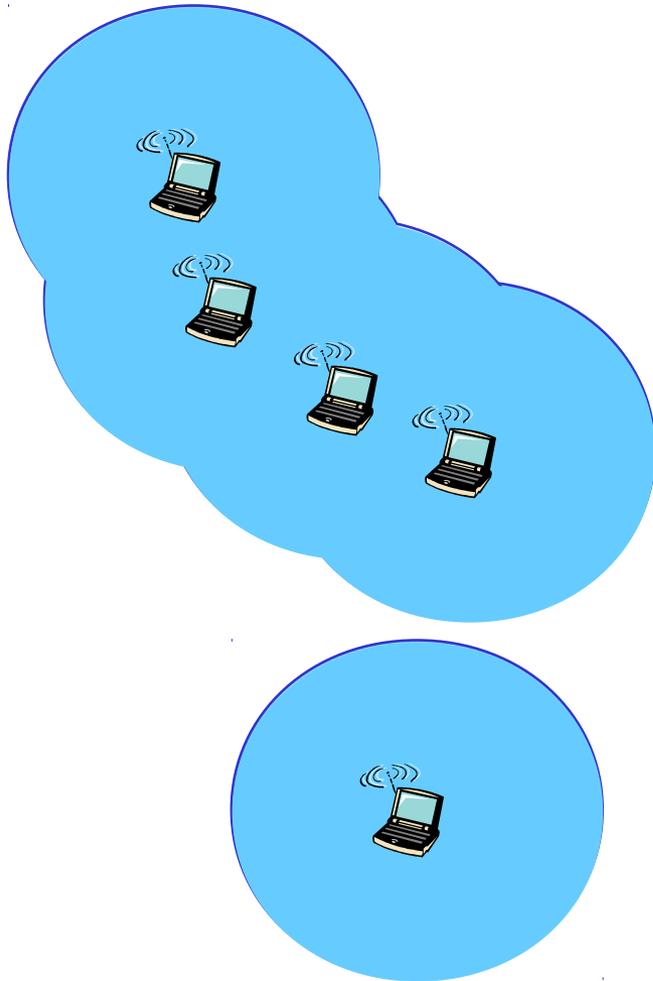
# Characteristics of selected wireless link standards



# Elements of a wireless network



# Elements of a wireless network



## Ad hoc mode

- no base stations
- nodes can only transmit to other nodes within link coverage
- nodes organize themselves into a network: route among themselves

# Taxonomy

- ❑ Single-hop, infrastructure-based
  - 802.11 infrastructure mode, 802.16 WiMAX
- ❑ Single-hop, infrastructure-less
  - 802.11 ad hoc mode, Bluetooth
- ❑ Multi-hop, infrastructure-based
  - Some wireless sensor networks
- ❑ Multi-hop, infrastructure-less
  - Mobile/vehicular ad hoc networks

# Wireless Link Characteristics

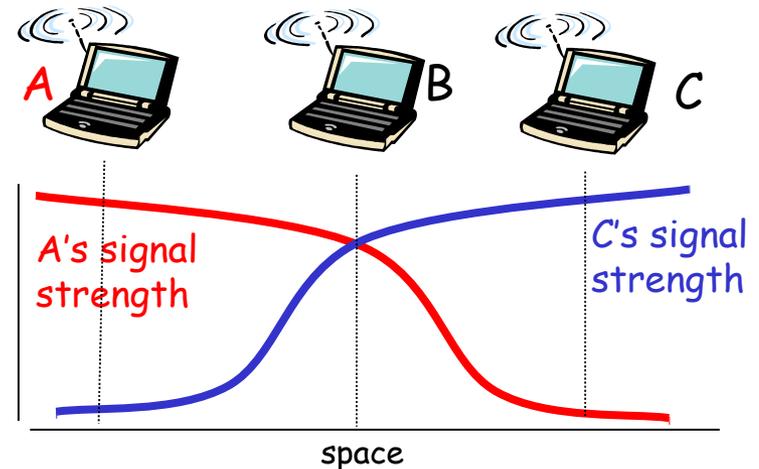
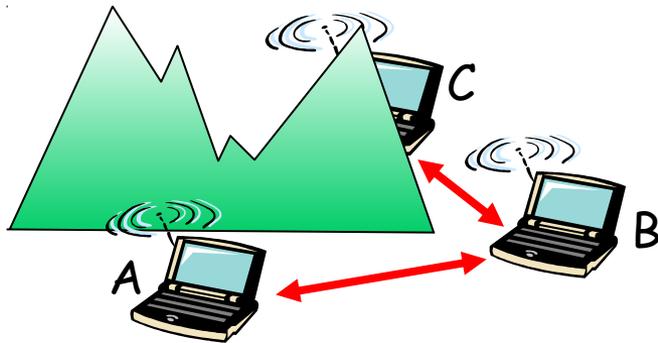
Differences from wired link ....

- **decreased signal strength:** radio signal attenuates as it propagates through matter (path loss)
- **interference from other sources:** standardized wireless network frequencies (e.g., 2.4 GHz for 802.11 {b,g}, Bluetooth, Zigbee) shared by other devices (e.g., US cordless phones, motors, microwaves)
- **multipath propagation:** radio signal reflects off objects ground, arriving at destination at slightly different times

... make communication across (even a point to point) wireless link much more "difficult"

# Wireless network characteristics

Multiple wireless senders and receivers create undetectable collision problems



Hidden terminal problem

- B, A hear each other
  - B, C hear each other
  - A, C can not hear each other
- means A, C unaware of their interference at B

Signal fading:

- B, A hear each other
  - B, C hear each other
  - A, C can not hear each other
- interfering at B

# Code Division Multiple Access (CDMA)

- ❑ Channel partitioning protocol
- ❑ used in several wireless broadcast channels (cellular, satellite, etc) standards
- ❑ unique "code" assigned to each user; i.e., code set partitioning
- ❑ all users share same frequency, but each user has own "chipping" sequence (i.e., code) to encode data
- ❑ *encoded signal* = (original data) X (chipping sequence)
  - $Z_{i,m} = d_i * c_m$
- ❑ *decoding*: inner-product of encoded signal and chipping sequence
  - $D_i = 1/M \sum Z_{i,m} * c_m$
- ❑ allows multiple users to "coexist" and transmit simultaneously with minimal interference (if codes are "orthogonal")

# Outline

## Introduction

### Wireless

- Wireless links, characteristics
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### Mobility

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## Summary

# IEEE 802.11 Wireless LAN

## □ 802.11b

- 2.4-5 GHz unlicensed radio spectrum
- up to 11 Mbps
- direct sequence spread spectrum (DSSS) in physical layer
  - all hosts use same chipping code
- widely deployed, using base stations

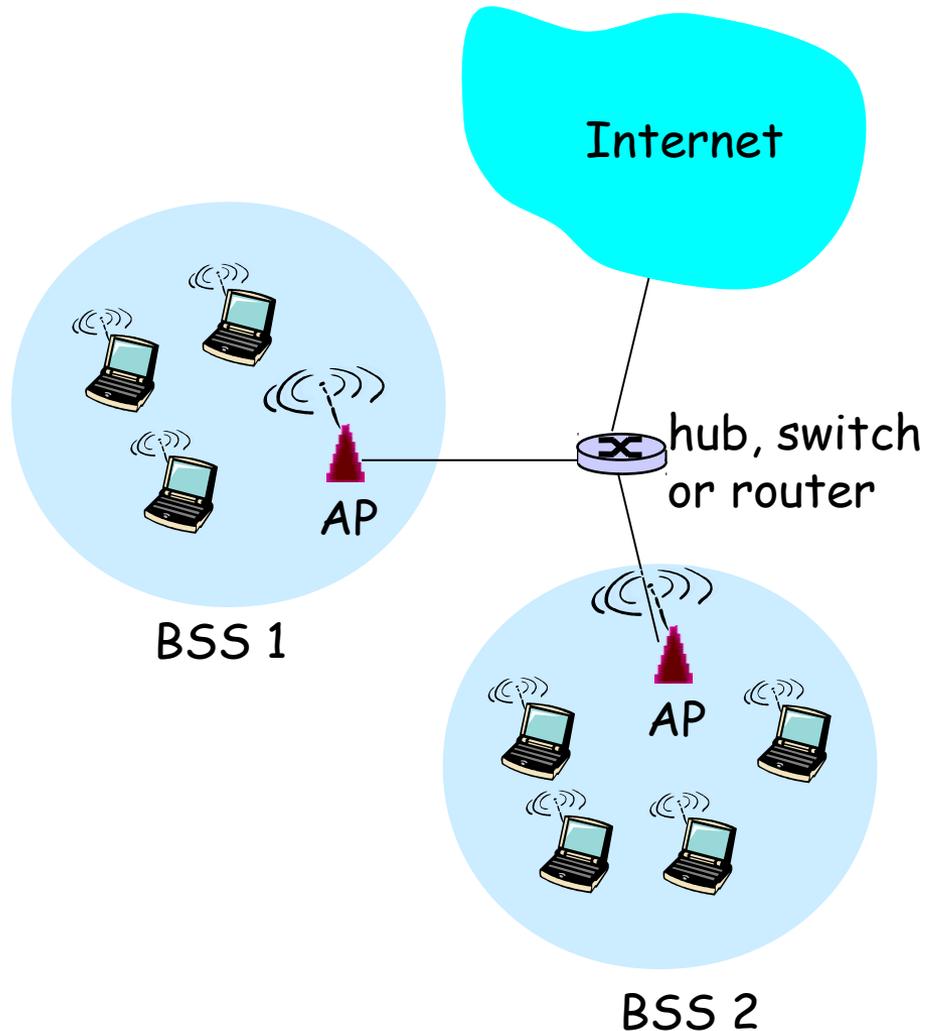
## □ 802.11a

- 5-6 GHz range
- up to 54 Mbps

## □ 802.11g

- 2.4-5 GHz range
- up to 54 Mbps
- All use *Carrier Sense Multiple Access with Collision Avoidance* for multiple access
- All have base-station and ad-hoc network versions

# 802.11 LAN architecture



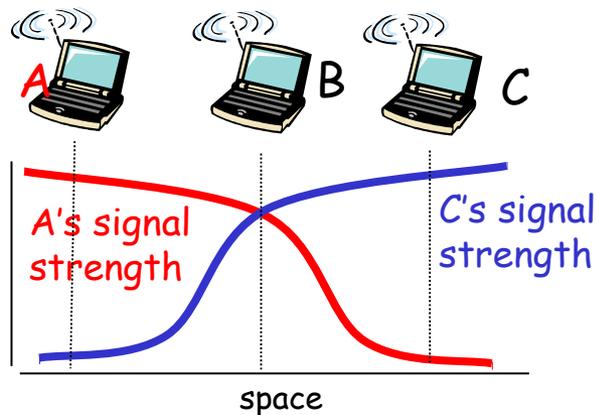
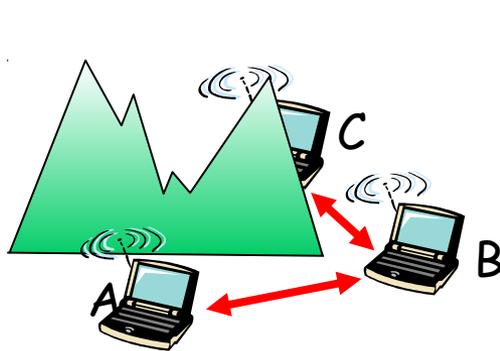
- wireless host communicates with base station
  - base station = access point (AP)
- Basic Service Set (BSS) (aka "cell") in infrastructure mode contains:
  - wireless hosts
  - access point (AP): base station
  - ad hoc mode: hosts only

# 802.11: Channels, association

- ❑ 802.11b: 2.4GHz-2.485GHz spectrum divided into 11 channels at different frequencies
  - AP admin chooses frequency for AP
  - interference possible: channel can be same as that chosen by neighboring AP!
- ❑ host: must *associate* with an AP
  - scans channels, listening for *beacon frames* containing AP's name (SSID) and MAC address
    - iwlist scan
  - selects AP to associate with
  - may perform authentication
    - iwconfig
  - will typically run DHCP to get IP address in AP's subnet
    - dhclient

# IEEE 802.11: multiple access

- ❑ avoid collisions: 2+ nodes transmitting at same time
- ❑ 802.11: CSMA - sense before transmitting
  - don't collide with ongoing transmission by other node
- ❑ 802.11: *no* collision detection!
  - difficult to receive (sense collisions) when transmitting due to weak received signals (fading)
  - can't sense all collisions in any case: hidden terminal, fading
  - goal: *avoid collisions*: CSMA/C(ollision)A(voidance)



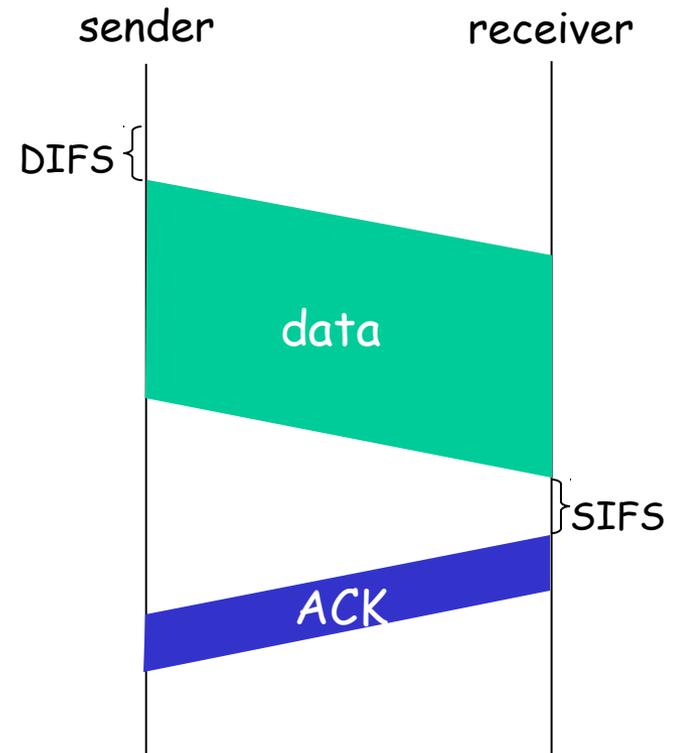
# IEEE 802.11 MAC Protocol: CSMA/CA

## 802.11 sender

- 1 if sense channel idle for **DIFS** then transmit entire frame (no CD)
- 2 if sense channel busy then start random backoff time  
timer counts down while channel idle  
transmit when timer expires  
if no ACK, increase random backoff interval, repeat 2

## 802.11 receiver

- if frame received OK  
return ACK after **SIFS** (ACK needed due to hidden terminal problem)

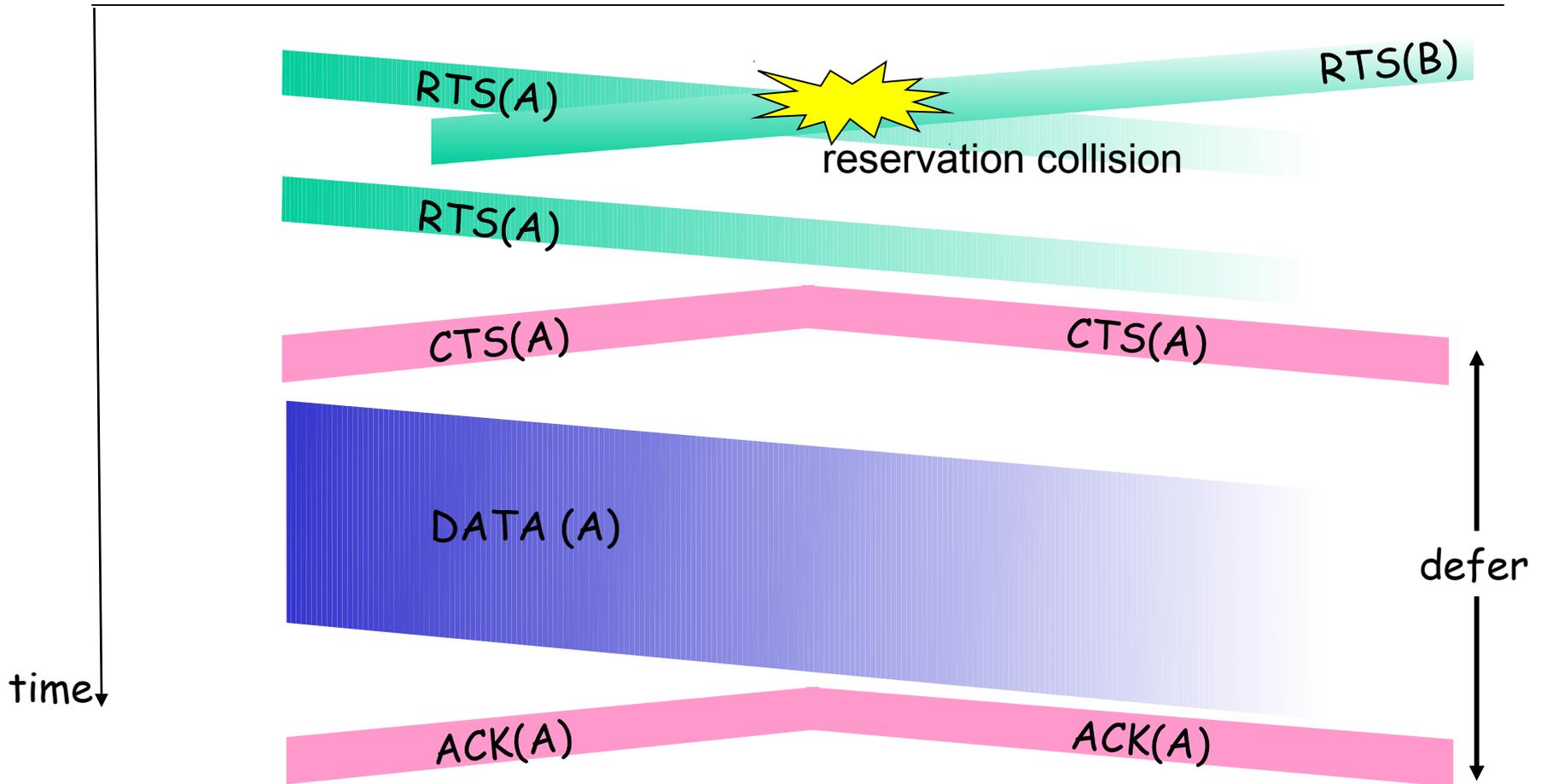


# Avoiding collisions (more)

- idea:* allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames
- sender first transmits *small* request-to-send (RTS) packets to BS using CSMA
    - RTSs may still collide with each other (but they’re short)
  - BS broadcasts clear-to-send CTS in response to RTS
  - RTS heard by all nodes
    - sender transmits data frame
    - other stations defer transmissions

Avoid data frame collisions completely  
using small reservation packets!

# Collision Avoidance with hidden terminals: Request-To-Send - Clear-To-Send exchange



# 802.11 frame: addressing



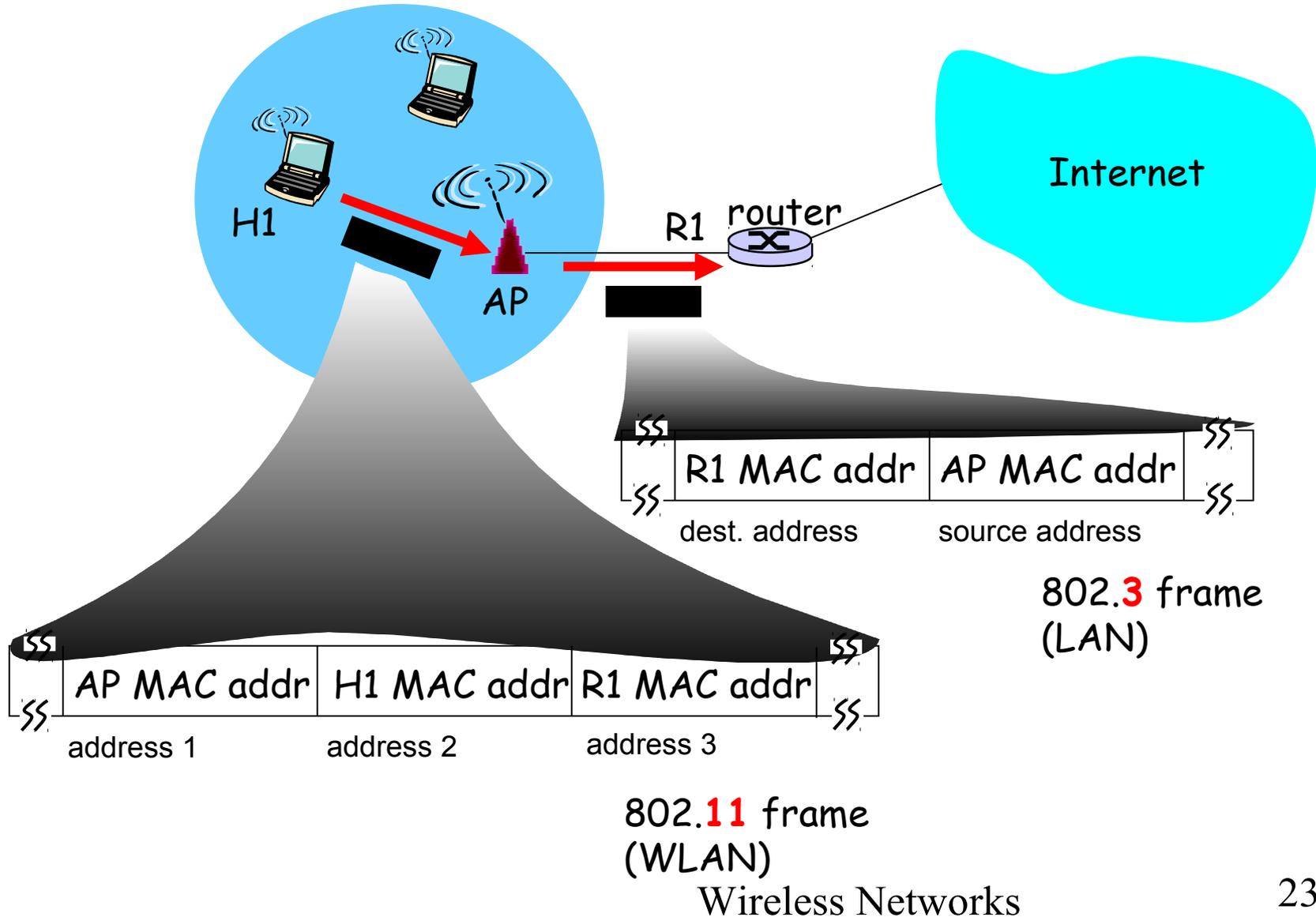
**Address 1:** MAC address of wireless host or AP to receive this frame

**Address 2:** MAC address of wireless host or AP transmitting this frame

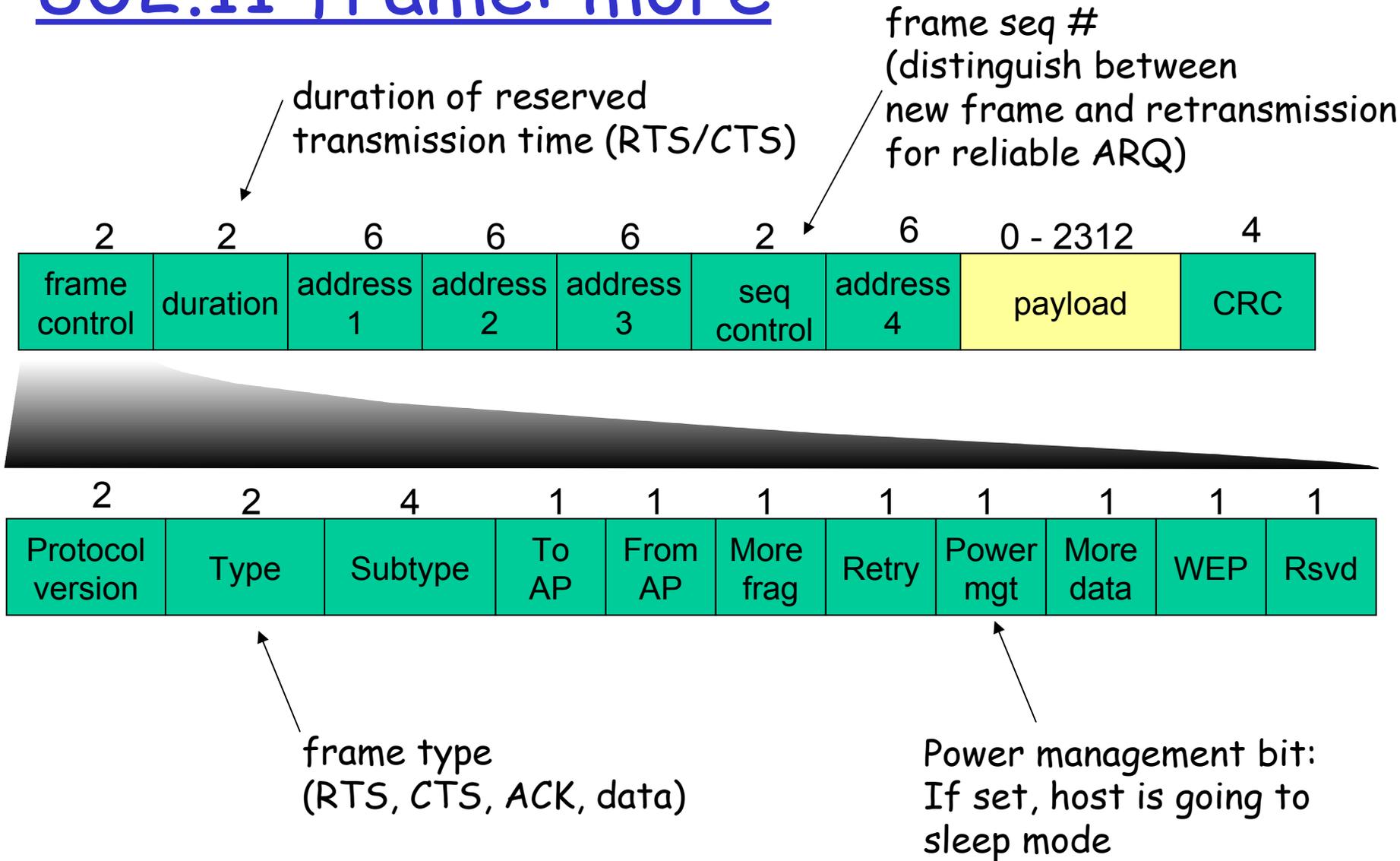
**Address 3:** MAC address of router interface to which AP is attached

**Address 4:** used only in ad hoc mode

# 802.11 frame: addressing

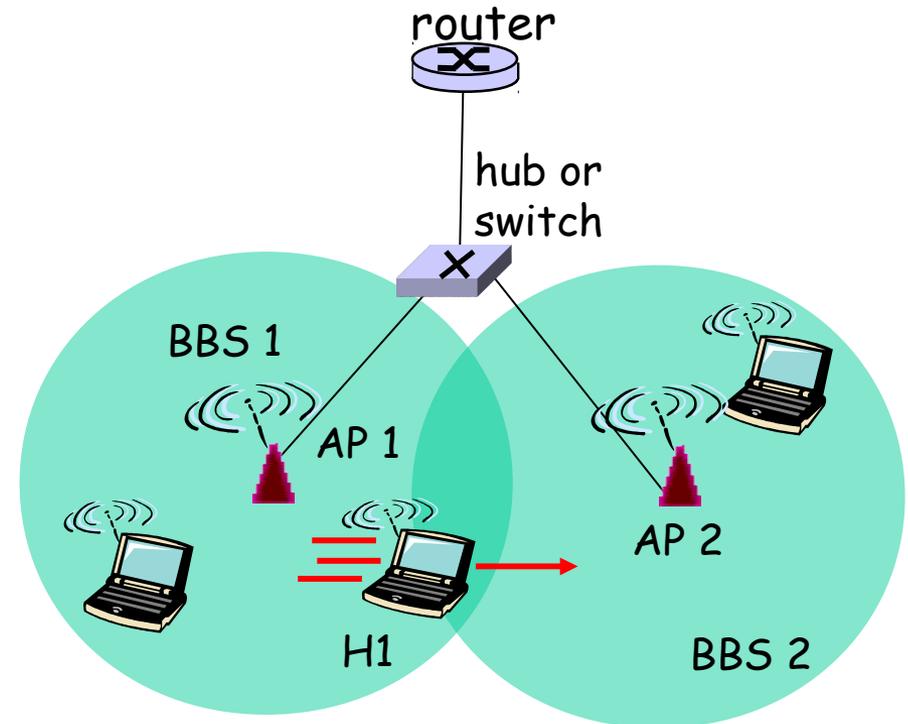


# 802.11 frame: more



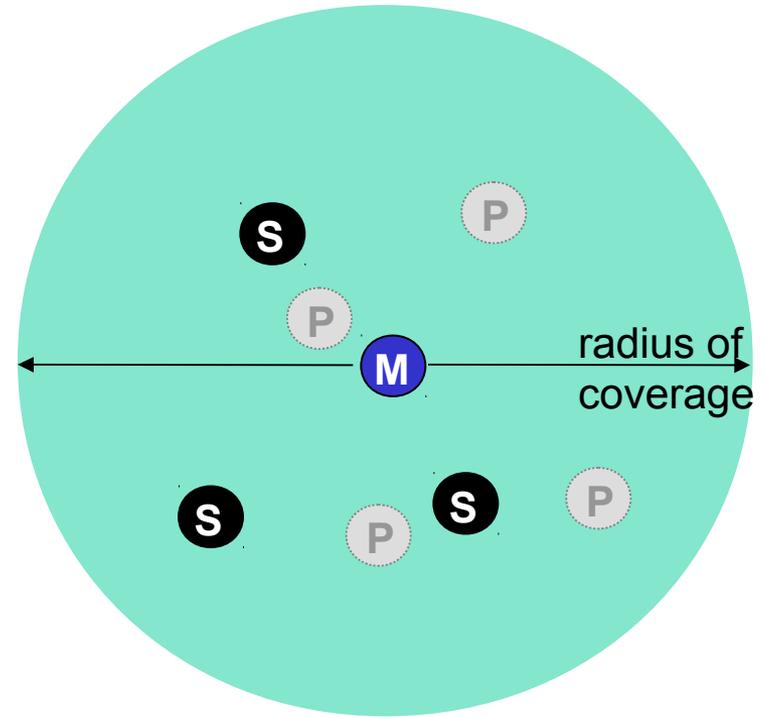
# 802.11: mobility within same subnet

- H1 remains in same IP subnet: IP address can remain same
- switch: which AP is associated with H1?
  - self-learning: switch will see frame from H1 and "remember" which switch port can be used to reach H1
  - AP1 & AP2 usually share SSID, so H1 handles the handoff easily



# 802.15: personal area network

- ❑ less than 10 m diameter
- ❑ replacement for cables (mouse, keyboard, headphones)
- ❑ ad hoc: no infrastructure
- ❑ master/slaves:
  - slaves request permission to send (to master)
  - master grants requests
- ❑ 802.15: evolved from Bluetooth specification
  - 2.4-2.5 GHz radio band
  - up to 721 kbps



- M** Master device
- S** Slave device
- P** Parked device (inactive)

# Outline

Introduction

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- Wireless links, characteristics
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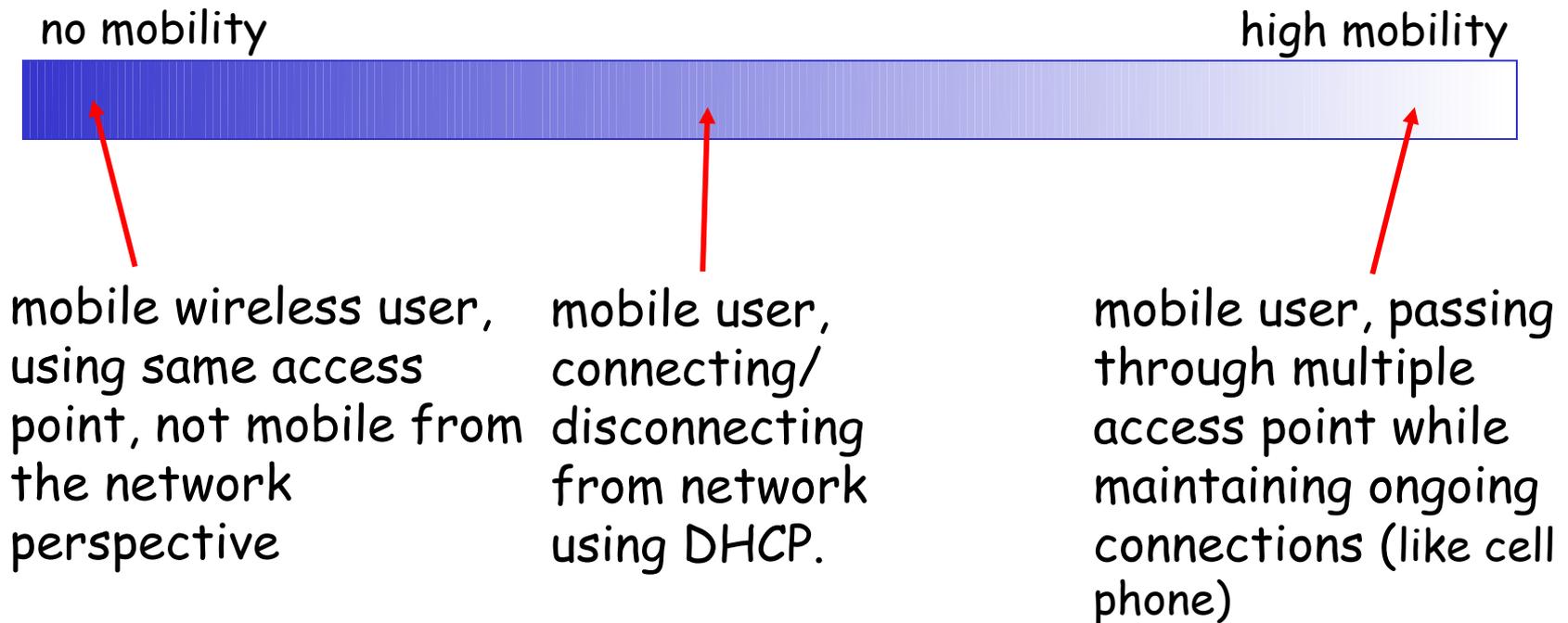
## Mobility

- Principles: addressing and routing to mobile users
- Mobile IP
- Mobility and higher-layer protocols

Summary

# What is mobility?

- spectrum of mobility, from the *network* perspective:

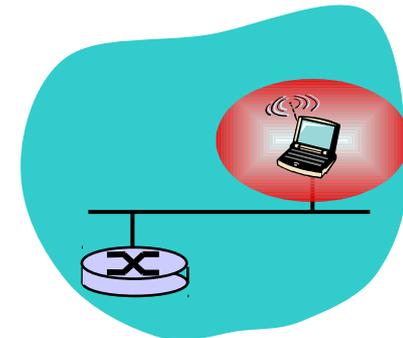
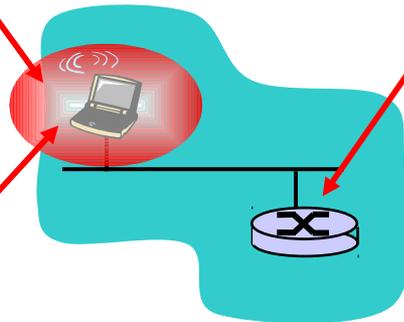


# Mobility: Vocabulary

**home network:** permanent  
"home" of mobile  
(e.g., 128.119.40/24)

**home agent:** entity that will  
perform mobility functions on  
behalf of mobile, when mobile  
is remote

**Permanent address:**  
address in home  
network, *can always* be  
used to reach mobile  
e.g., 128.119.40.186

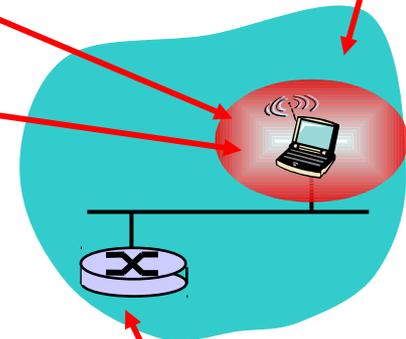
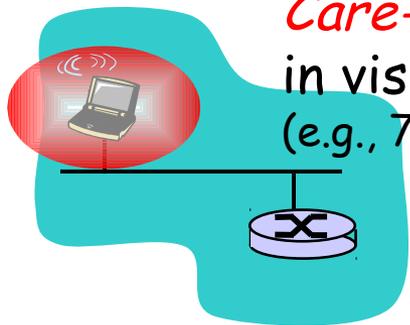


# Mobility: more vocabulary

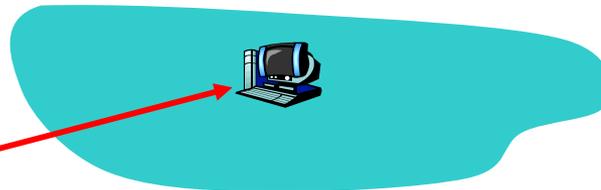
*Permanent address:* remains constant (e.g., 128.119.40.186)

*visited network:* network in which mobile currently resides (e.g., 79.129.13/24)

*Care-of-address:* address in visited network. (e.g., 79.129.13.2)



*correspondent:* wants to communicate with mobile

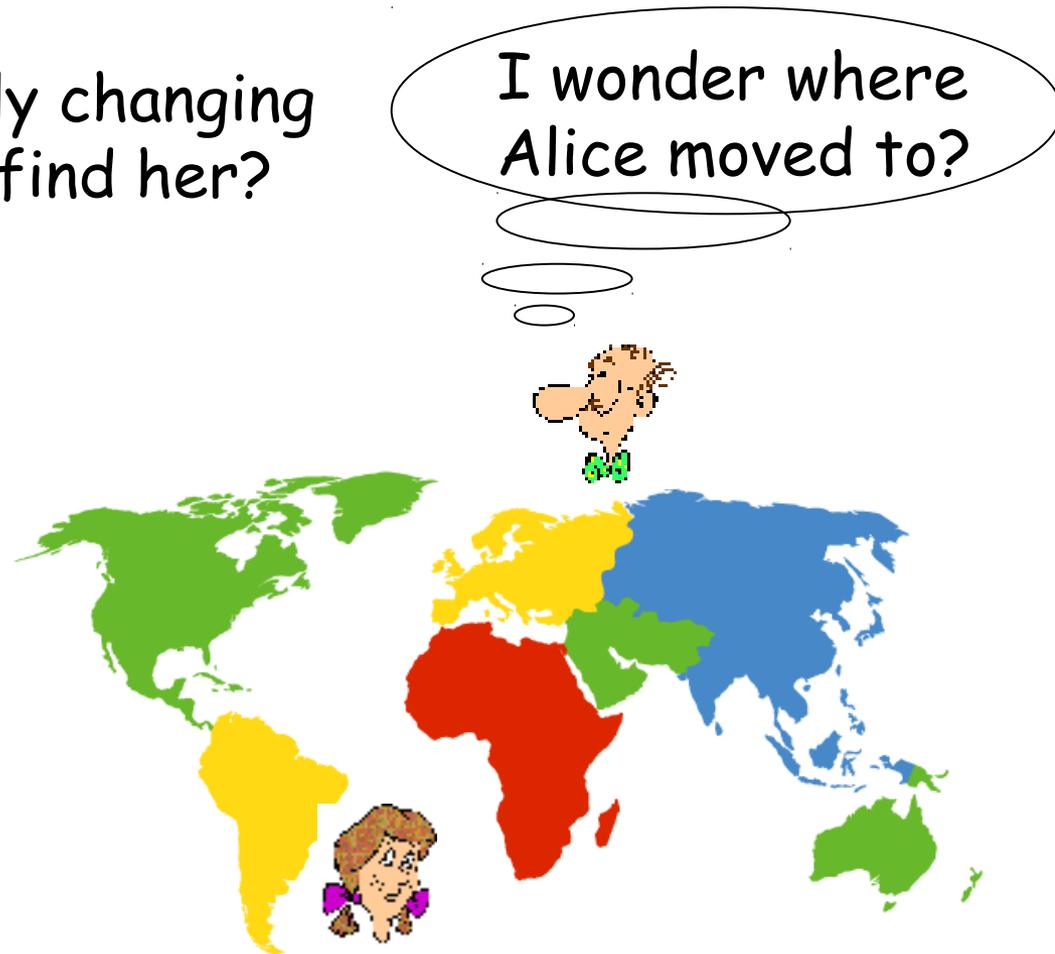


*foreign agent:* entity in visited network that performs mobility functions on behalf of mobile.

# How do you contact a mobile friend:

Consider friend frequently changing addresses, how do you find her?

- search all phone books?
- call her parents?
- expect her to let you know where he/she is?



# Mobility: approaches

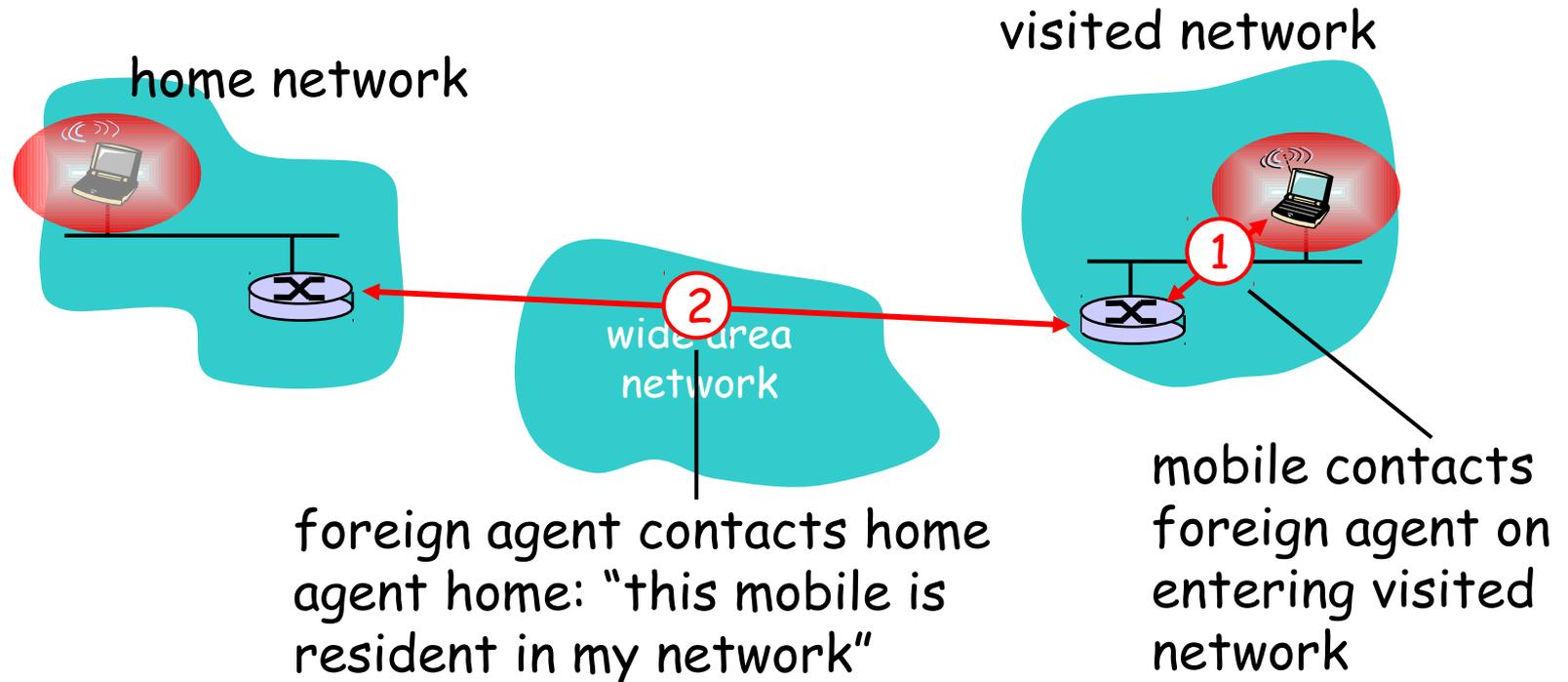
- *Let routing handle it:* routers advertise permanent address of mobile-nodes-in-residence via usual routing table exchange.
  - routing tables indicate where each mobile located
  - no changes to end-systems
- *Let end-systems handle it:*
  - *indirect routing:* communication from correspondent to mobile goes through home agent, then forwarded to remote
  - *direct routing:* correspondent gets foreign address of mobile, sends directly to mobile

# Mobility: approaches

- ❑ *Let routing handle it:* routers advertise permanent address of mobile, residence via usual routing table entries
  - routing table entries for where each mobile located
  - no changes to end systems
- ❑ *let end-systems handle it:*
  - *indirect routing:* communication from correspondent to mobile goes through home agent, then forwarded to remote
  - *direct routing:* correspondent gets foreign address of mobile, sends directly to mobile

not  
scalable  
to millions of  
mobiles

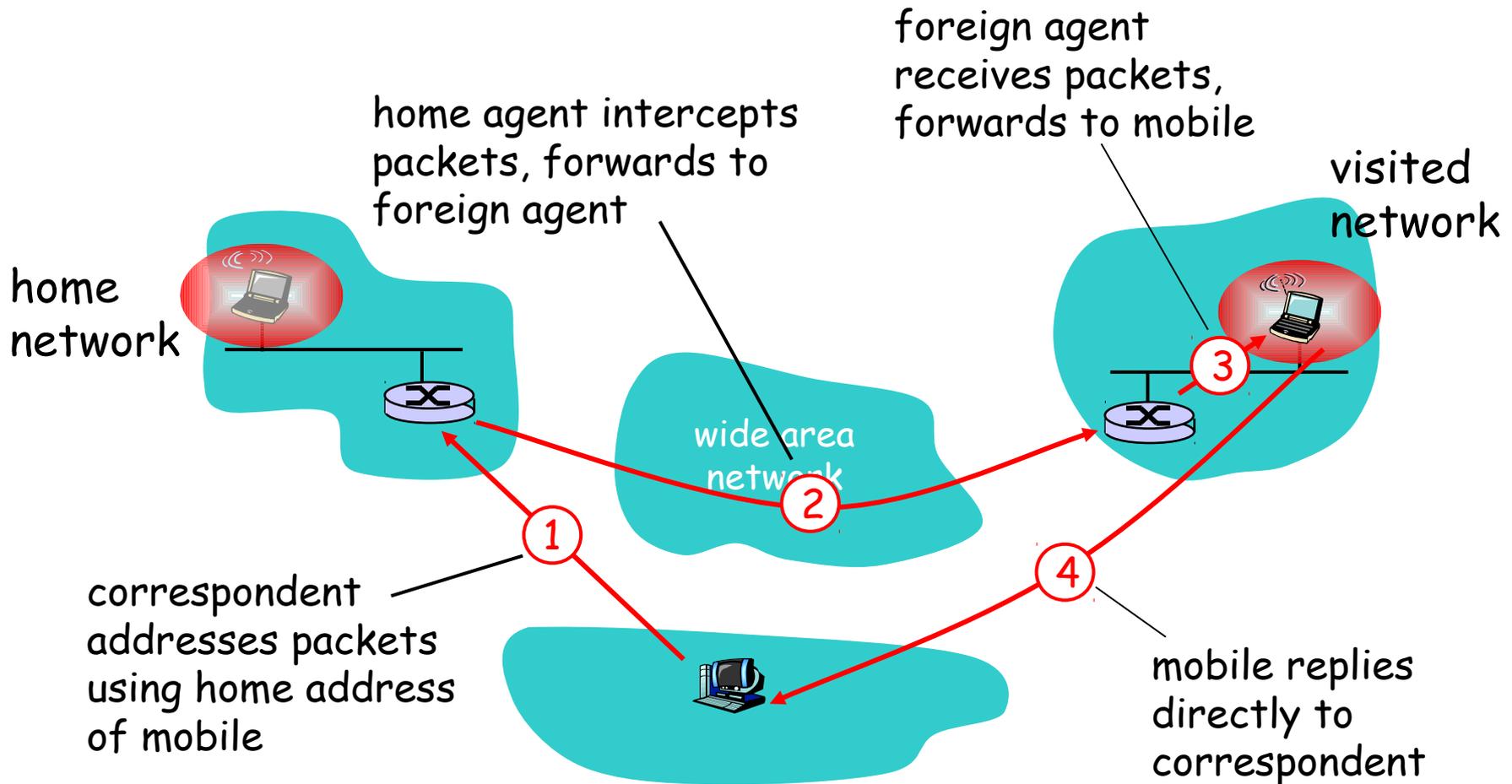
# Mobility: registration



End result:

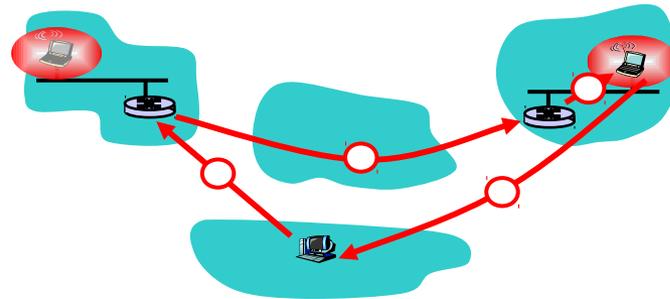
- ❑ Foreign agent knows about mobile
- ❑ Home agent knows location of mobile

# Mobility via Indirect Routing



# Indirect Routing: comments

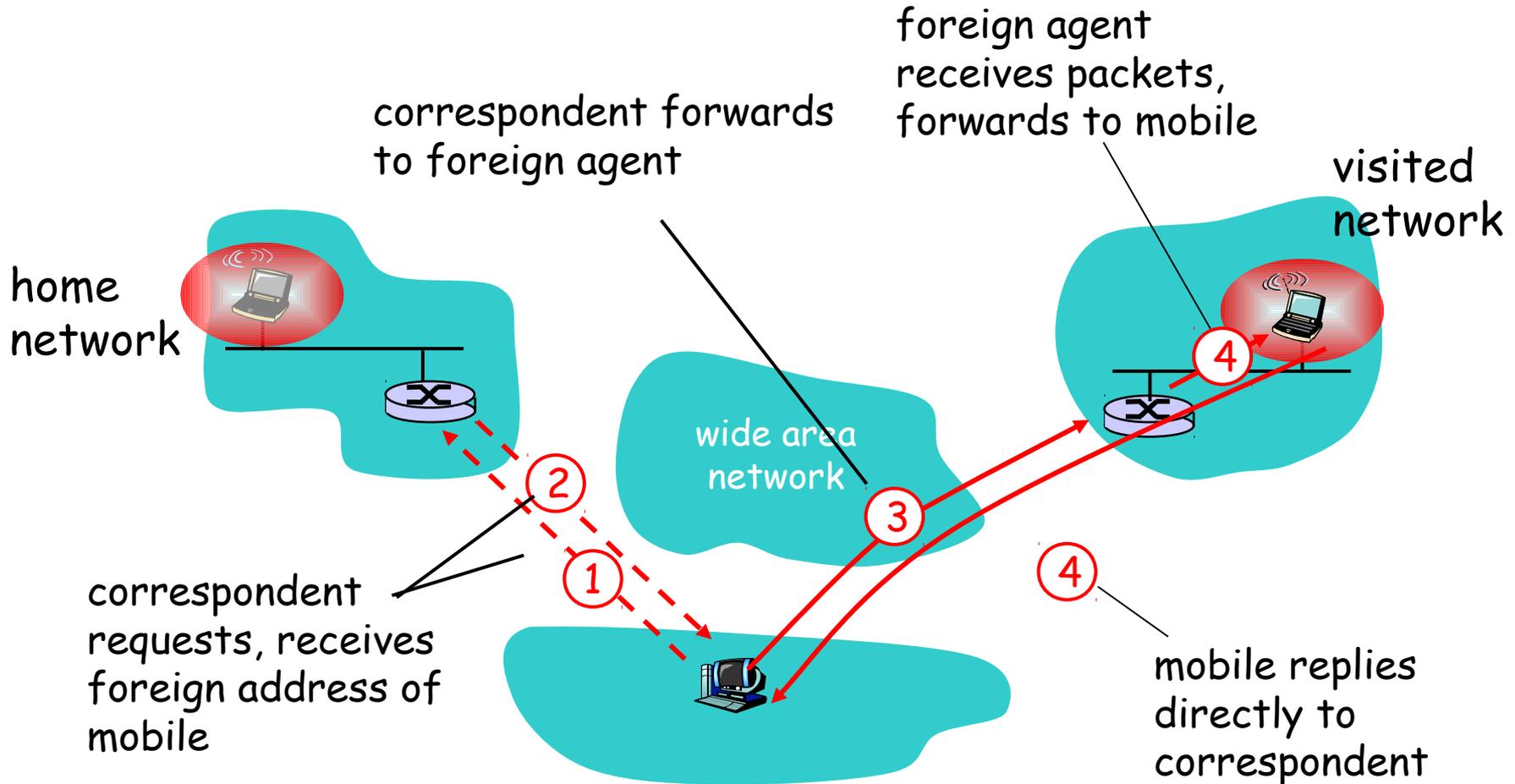
- Mobile uses two addresses:
  - permanent address: used by correspondent (hence mobile location is *transparent* to correspondent)
  - care-of-address: used by home agent to forward datagrams to mobile
- foreign agent functions may be done by mobile itself
- triangle routing: correspondent-home-network-mobile
  - inefficient when correspondent, mobile are in same network



# Indirect Routing: moving between networks

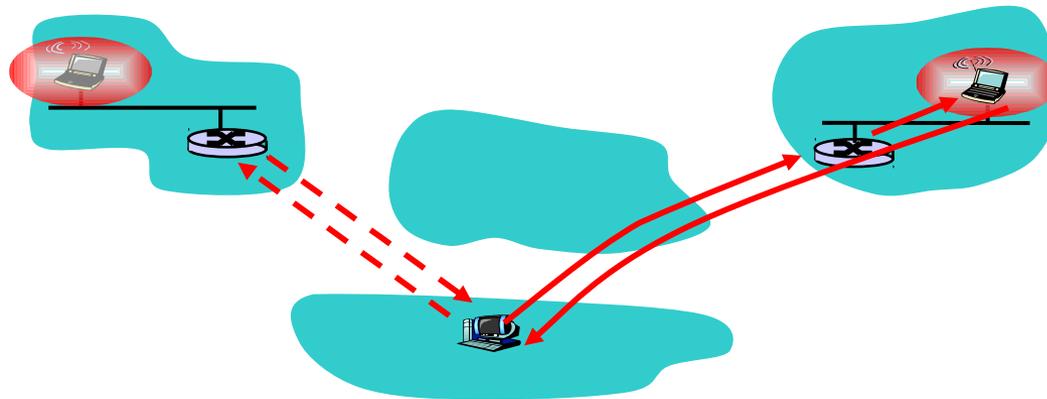
- suppose mobile user moves to another network
  - registers with new foreign agent
  - new foreign agent registers with home agent
  - home agent update care-of-address for mobile
  - packets continue to be forwarded to mobile (but with new care-of-address)
- mobility, changing foreign networks transparent: *on going connections can be maintained!*

# Mobility via Direct Routing



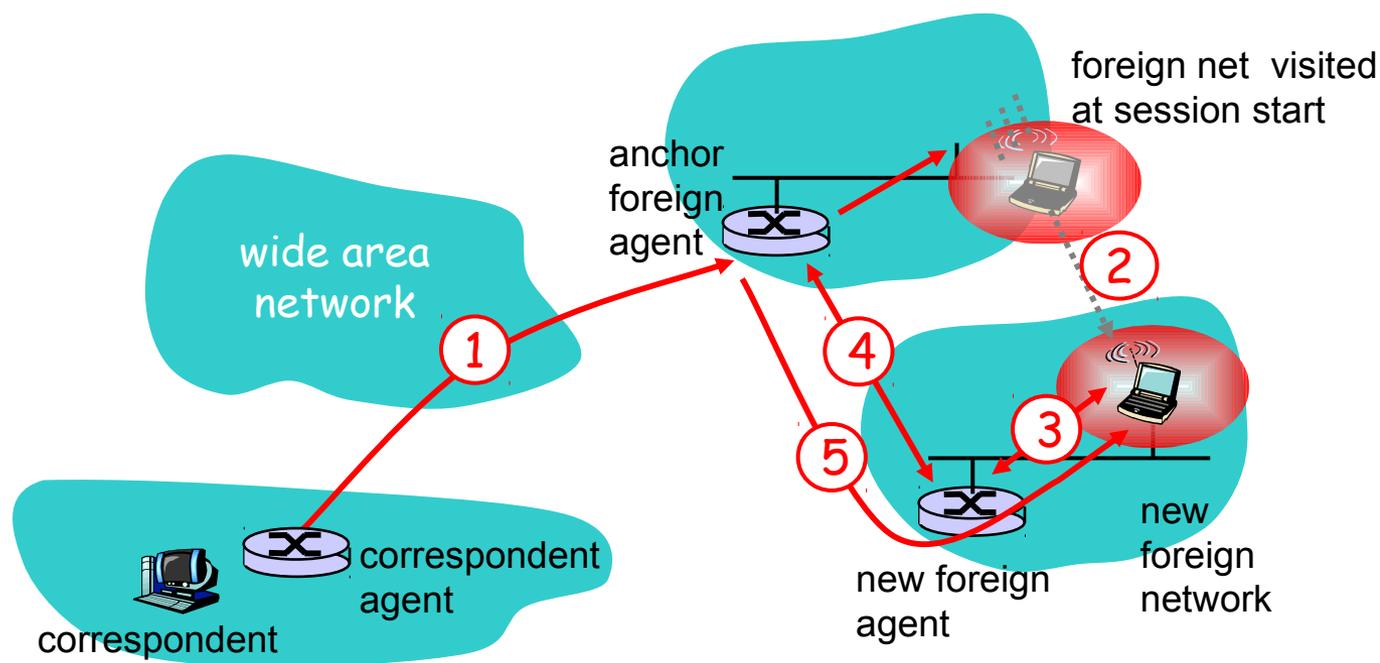
# Mobility via Direct Routing: comments

- overcome triangle routing problem
- **non-transparent to correspondent:**  
correspondent must get care-of-address from home agent
  - what if mobile changes visited network?



# Accommodating mobility with direct routing

- ❑ anchor foreign agent: FA in first visited network
- ❑ data always routed first to anchor FA
- ❑ when mobile moves: new FA arranges to have data forwarded from old FA (chaining)



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### Mobility

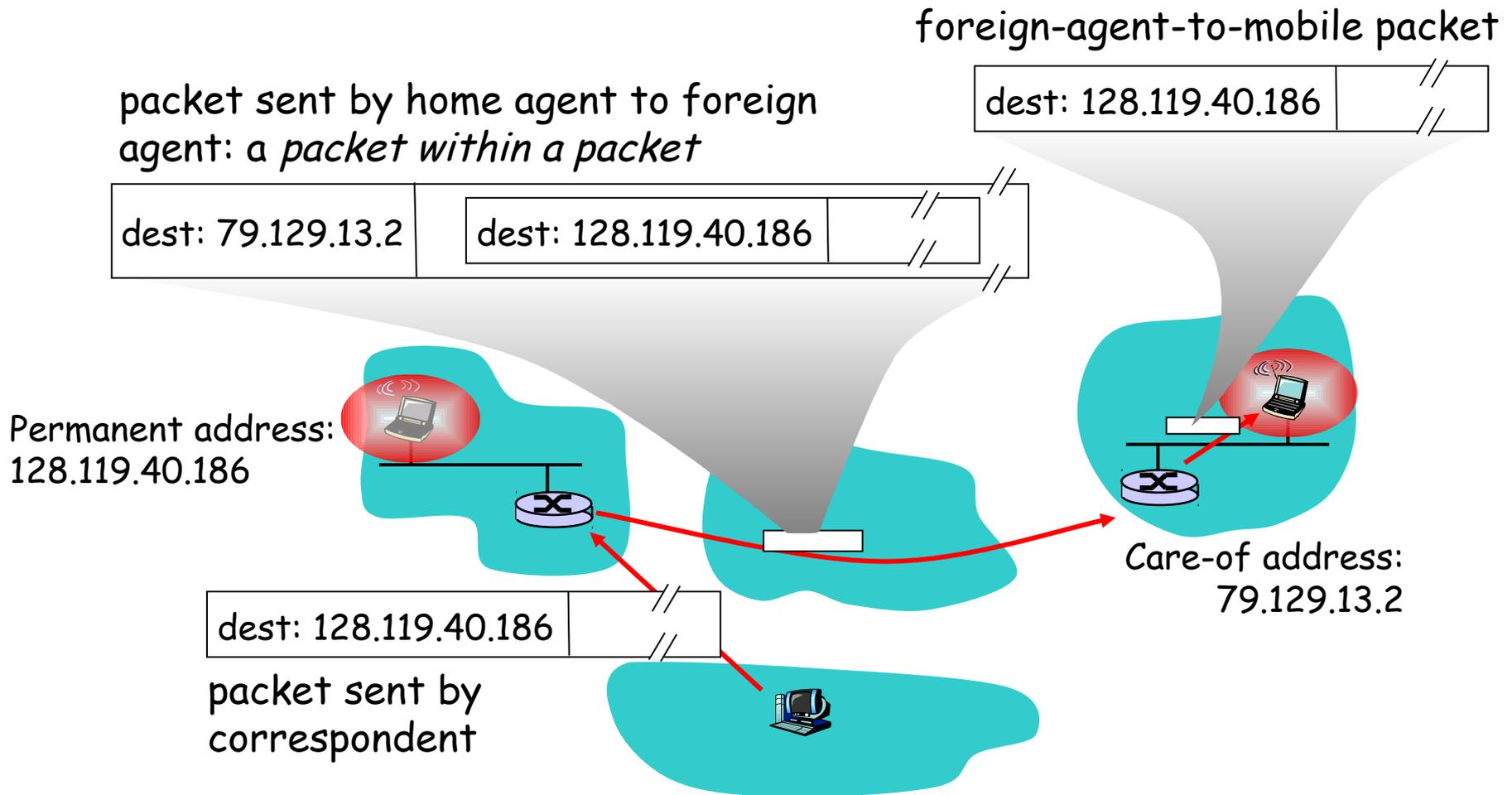
- Principles: addressing and routing to mobile users
- **Mobile IP**
- Mobility and higher-layer protocols

## Summary

# Mobile IP

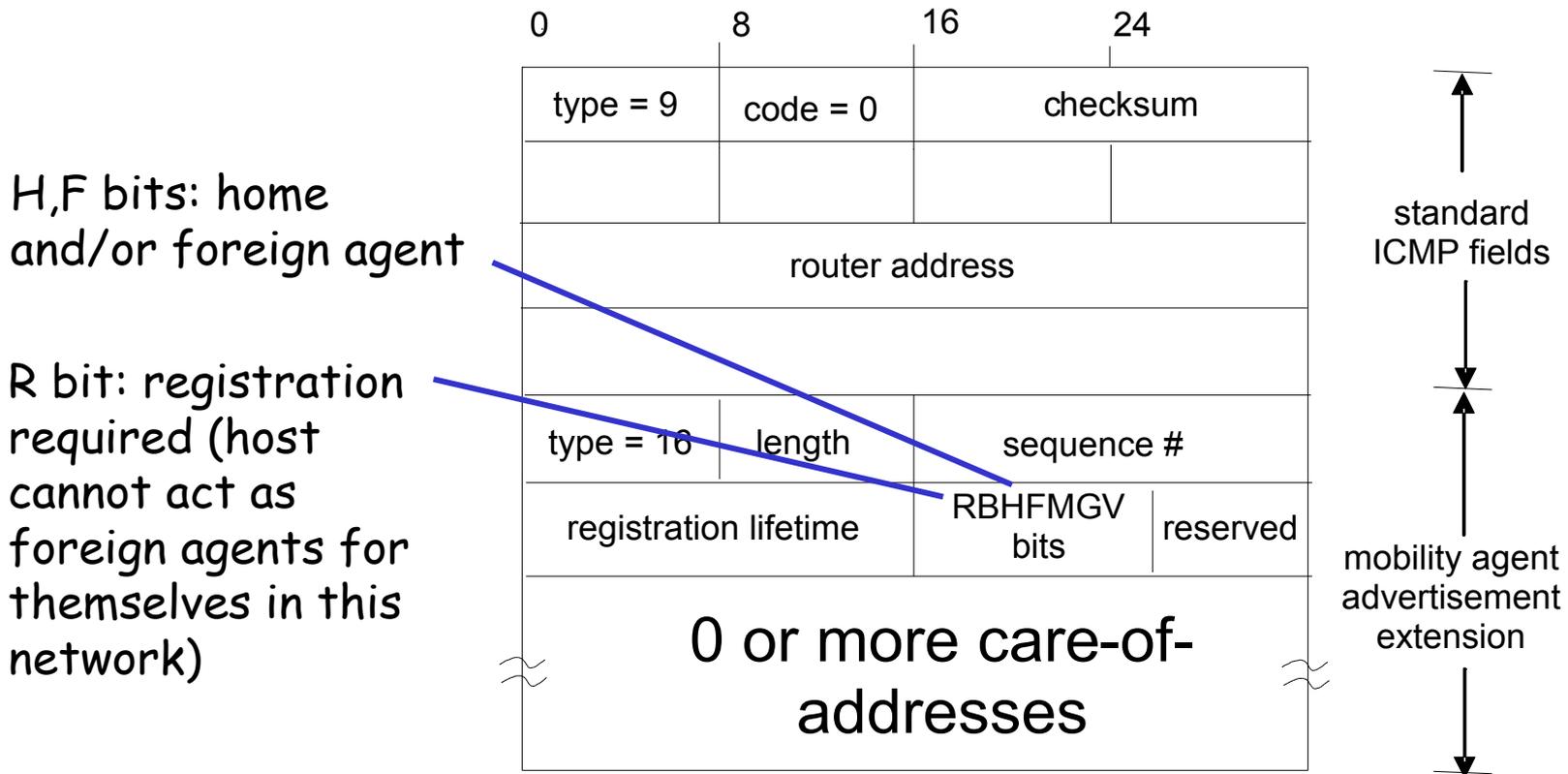
- ❑ RFC 3344 for IPv4
- ❑ has many features we've seen:
  - home agents, foreign agents, foreign-agent registration, care-of-addresses, encapsulation (packet-within-a-packet)
- ❑ three components to standard:
  - indirect routing of datagrams
  - agent discovery
  - registration with home agent
- ❑ major security considerations

# Mobile IP: indirect routing

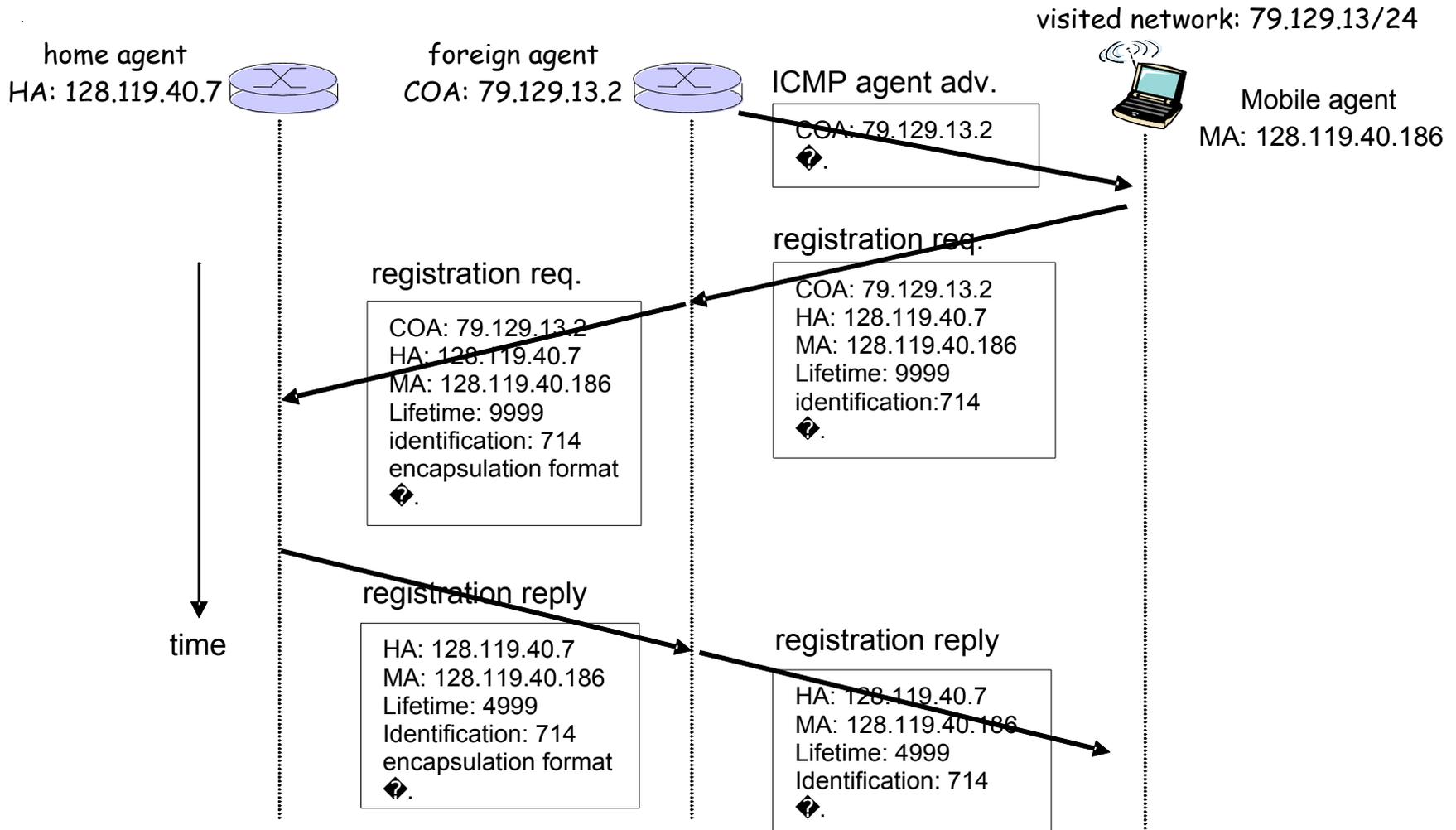


# Mobile IP: agent discovery

- **agent advertisement:** foreign/home agents advertise service by broadcasting ICMP control messages (typefield = 9, router discovery)



# Mobile IP: registration example



## Wireless, mobility: impact on higher layer protocols

- ❑ logically, impact *should* be minimal ...
  - best effort service model remains unchanged
  - TCP and UDP can (and do) run over wireless, mobile
- ❑ ... but performance-wise:
  - packet loss/delay due to bit-errors (discarded packets, delays for link-layer retransmissions), and handoff
  - TCP interprets loss as congestion, will decrease congestion window un-necessarily
  - delay impairments for real-time traffic
  - limited bandwidth of wireless links

# Summary

## Wireless

- wireless links:
  - capacity, distance
  - channel impairments
  - CDMA
- IEEE 802.11 (“wi-fi”)
  - CSMA/CA reflects wireless channel characteristics

## Mobility

- principles: addressing, routing to mobile users
  - home, visited networks
  - direct, indirect routing
  - care-of-addresses
- case study
  - mobile IP
- impact on higher-layer protocols